### An Axiomatic Framework for Multi-channel Attribution in Online Advertising

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### **Optimization Cycle in Online Advertising**



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A lot of academic effort and progress on media mix optimization ... relatively less progress on the measurement and attribution steps

### Multi-Channel Attribution in Online Advertising



Question: How to give credit to all channels contributing to a conversion?

#### Common current attribution methods



Source: Beyond Last Touch: Understanding Campaign Effectiveness, 2013 Quantcast

## Goals for this work

- Framework for attributing credit to channels
- Evaluate heuristic procedures
- Address limitations of current heuristics and propose prescriptions

#### **Related work**

#### Heuristics for multi-touch attribution

- Shapley value: [Shao and Li, 2011], [Dalessandro et al., 2012], [Berman, 2013], [Li and Kannan, 2014]
- Incremental impact: [Abhishek et al., 2012], [Anderl et al., 2014]

#### **Empirical work**

• [Kireyev et al., 2013], [Lewis and Rao, 2014]

#### Other related work

• [Jordan et al., 2011], [Archak et al., 2010]





Two possible effects of an action:



Two possible effects of an action: Network effect



Two possible effects of an action: Network effect and Funnel effect

# Approach

"Observational" inference

- Each channel c shows the ad to  $\beta_c$  fraction of customers
- Observe one-step browsing behavior

Attribution

- Payment per impression
- Payment independent of the customer state or browsing behavior

## Model

Markov chain consisting two state variables  $\sigma=(c,s)$ 

- $\bullet \ c = {\rm identity} \ {\rm of} \ {\rm the} \ {\rm channel, \ e.g.} \ {\rm particular} \ {\rm website}$
- s =latent state of the visitor

Will ignore latent state in the talk. All results go through!

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Transition probability

- $P_{c,c'}^{(1)}$  = probability  $c \rightarrow c'$  when a visitor exposed to an ad on channel c
- $P_{c,c'}^{(0)} =$  probability  $c \rightarrow c'$  when a visitor not exposed to an ad on channel c
- $\bullet$  Transition when  $\beta$  fraction see an ad

$$P^{\beta} = \operatorname{diag}(\beta)P^{(1)} + \operatorname{diag}(1-\beta)P^{(0)}$$

# Model (contd)

- Two absorbing states
  - +  $c^{\ast}$  = conversion. Conversion only on advertising.
  - q = quit
- $\pi_c^* = \text{conversion probability from channel } c$

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- Two absorbing states
  - $c^* = \text{conversion}$ . Conversion only on advertising.
  - q = quit
- $\pi_c^* = \text{conversion probability from channel } c$
- $\lambda_c =$  external traffic to channel c and  $\mu_c^\beta =$  effective traffic at c

$$\mu^{\beta} = \lambda + (P^{\beta})^{\top} \mu^{\beta}$$

•  $h_c^{\beta} = \text{conversion probability from } c$ 

$$h^\beta = P^\beta h^\beta + \pi^*$$

#### Payment per impression

Channel c is paid  $z_c$  per impression

- Payment not a function of the latent state
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Desirable properties

- Channel rationality:  $z \ge 0$
- Advertiser rationality: global budget balance

Total value generated

 $\geq (\mu^{\beta})^{\top} \operatorname{diag}(\beta) z$ 

Total payments

• Fairness: hard to quantity!

Pay each channel one unit for each conversion

Not quite allowed in our formulation but admits equivalent representation through payment per impression

Trivially satisfies global budget balance. But, very unfair!

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Incremental value of the advertising action

- Advertise:  $P^{(1)}h^{\beta} + \pi^*$
- Do not advertise:  $P^{(0)}h^{\beta}$
- Incremental value:  $\underbrace{(P^{(1)} P^{(0)})}_{:= \Lambda} h^{\beta} + \pi^*$

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Incremental value heuristic:  $z = \Delta h^{\beta} + \pi^*$ . Pays for both

- immediate conversion:  $\pi^*$
- network effect:  $\Delta h^{\beta}$

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**Example**:  $z_c \equiv 1$  fair! But  $\sum_c \mu_c^\beta \beta_c z_c = n \gg \sum_c \lambda_c h_c^\beta = 1!$ 



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## IVH: generically overbudget

Proposition. Let  $\delta = \mu^T \operatorname{diag}(\beta) \Delta h(\beta)$ . Then  $\lambda^\top h^\beta - (\mu^\beta)^\top \operatorname{diag}(\beta) z = -\delta$ 

 $\delta>0\approx$  when advertising generates downstream indirect conversions

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 $\delta>0\approx$  when advertising generates downstream indirect conversions

Over payment: Every conversion gets counted several times

- pay for both for direct conversion, and
- eventual conversion

### IVH with past payments

Need to account for past payments!

Proposal:  $z_c = (\Delta h)_c + \pi_c^* - \mathbb{E}[\text{past payments}|c]$ 

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- Channel a: cannot be compensated for the  $1 \epsilon$  fraction it converts!

Requirements

- discriminate between outcomes when allocating payments
- local budget balance must ensure global budget balance

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Allocate cost  $x_{a,b}$  along the arc (a,b)

• Axiom 1: Cost proportional to outcome value:  $x_{a,b} \propto h_b^\beta$ 

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• Axiom 3: Backward payment conservation:  $p_c = \sum_b \mathbb{P}(b \mid c) x_{b,c}$ 

#### Outcome based accounting (contd)

Axioms 1+2

$$x_{a,b} = (p_a + \beta_a z_a) \cdot \frac{h_b^\beta}{\sum_c P_{a,c}^\beta h_c^\beta + \pi_a^*} = (p_a + \beta_a z_a) \cdot \frac{h_b^\beta}{h_a^\beta}$$

Axiom 3

$$p_c = \sum_{b} \mathbb{P}(b \mid c) x_{b,c}$$
$$= \sum_{b} \left( \frac{\mu_b P_{bc}^{\beta}}{\mu_c} \right) \cdot \frac{h_c}{h_b} \cdot (p_b + \beta_b z_b)$$

**Definition.** (Local budget balance)

$$p_c + \beta_c z_c \le h_c^\beta$$





#### Past payments

• 
$$p_a = 0$$
  
•  $p_b = z_a \cdot \frac{h_b}{h_a} = z_a \cdot \frac{\epsilon}{1 - \epsilon + \epsilon^2}$ 



#### Past payments

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Budget balance conditions

• 
$$z_a \le h_a = 1 - \epsilon + \epsilon^2$$
  
•  $z_b + \frac{\epsilon z_a}{1 - \epsilon + \epsilon^2} \le h_b = \epsilon$ 



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$$z_a \le h_a = 1 - \epsilon + \epsilon^2$$
  
•  $z_b + \frac{\epsilon z_a}{1 - \epsilon + \epsilon^2} \le h_b = \epsilon$ 

$$z_a = 1 - 2\epsilon + 2\epsilon^2 - \epsilon^3$$
 and  $z_b = \epsilon^2$  feasible!



Past payments

$$p_1 = 0$$
  

$$p_k = (p_{k-1} + z_{k-1}) \cdot \frac{h_k}{h_{k-1}} = p_{k-1} + z_{k-1} = \sum_{j < k} z_j$$



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Local budget balance

$$p_k + z_k \le h_k = 1, \forall k \quad \Leftrightarrow \quad \sum_{k=1}^n z_k \le 1$$

Global budget balance!

### Local balance implies global balance!

#### Theorem. Suppose

$$p_c + eta_c z_c \ \leq \ h_c^eta, \ \ orall c$$
 (Local budget balance)

Then

$$(\mu^{\beta})^{\top} \operatorname{diag}(\beta) z \leq \lambda^{\top} h^{\beta}.$$
 (Global budget balance)

#### Proof.

- $\bullet$  Use definition of  $p_c$
- Sum all local budget balance constraint
- Careful accounting of slacks

#### Feasible per impression payments z

#### Channel rationality: $z \ge 0$

#### Advertiser rationality: global budget balance implied by

- Outcome based cost allocation
- Local budget balance

#### Fairness

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#### Fairness

• IVH with payments: Value  $V_c$  at channel c satisfies

$$V = P^{\beta}V + \pi^* - \mathbf{diag}(\beta)z$$

Reduces to  $h^\beta$  when  $z\equiv 0$ 

 $z \leq \Delta V + \pi^*$  [No lift  $\implies$  No payment]

Does **not** identify z uniquely! Introduce an objective function.

## Maximize payments

#### Example



Max payment any payment satisfying global budget constraint  $\sum_k z_k \leq 1$ 

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Very fragile!

•  $\pi_1^* = \epsilon$ : unique optimal solution:  $z_1^* = h_1$  and  $z_j^* = 0$  for j > 1

### Nash bargaining solution

Let  $\mathcal{Z}$  denote a set of feasible payments for a given network.

$$\mathcal{Z} = \left\{ z \mid \begin{array}{l} V = P^{\beta}V + \pi^* - \operatorname{diag}(\beta)z \\ z \leq \Delta V + \pi^* \\ p + \operatorname{diag}(\beta)z \leq h^{\beta} \\ p_c = \sum_b \left(\frac{\mu_b P_{bc}^{\beta}}{\mu_c}\right) \cdot \frac{h_c}{h_b} \cdot (p_b + \beta_b z_b) \end{array} \right\}$$

## Nash bargaining solution

Let  $\ensuremath{\mathcal{Z}}$  denote a set of feasible payments for a given network.

Let f denote a function that maps  $\mathcal Z$  to an element  $z\in \mathcal Z$  that satisfies

- Pareto efficiency
- Symmetry
- Invariance to scaling payoffs
- Independence of Irrelevant Alternatives

Then

$$f(\mathcal{Z}) = \max_{z} \sum_{i=1}^{n} \ln(z_i)$$
  
s.t.  $z \in \mathcal{Z}$ 



Nash bargaining solution:  $\mu_i \beta_i z_i = \mu_i h_i / n$ 

	$z_1$	$z_2$	$z_3$	$z_4$	Total
Max	0.0000	0.0000	0.0000	0.5000	0.0625
Nash	0.0156	0.0313	0.0625	0.1250	0.0625
Nash + IVH	0.0156	0.0312	0.0625	0.1250	0.0625

#### Sensitivity analysis: Path



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Change  $\pi_1^* = 0.1$ 

	$z_1$	$z_2$	$z_3$	$z_4$	Total
Max	0.225	0.0	0.0	0.0	0.225
Nash	0.0406	0.0313	0.0625	0.125	0.0875
Nash + IVH	0.0406	0.0313	0.0625	0.125	0.0875

#### Sensitivity analysis: Path



Nash bargaining solution:  $\mu_i\beta_i z_i = \mu_i h_i/n$ 

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Nash + IVH	0.0156	0.0312	0.0625	0.1250	0.0625

Set  $P^{(0)}(3,4) = P^{(1)}(3,4) = 0.5$ 

	$z_1$	$z_2$	$z_3$	$z_4$	Total
Max	0.0000	0.0000	0.0000	0.5000	0.0625
Nash	0.0156	0.0313	0.0625	0.1250	0.0625
Nash + IVH	0.0234	0.0469	0.0000	0.1250	0.0625

#### Sensitivity analysis: Cycle



	$z_1$	$z_2$	$z_3$	$z_4$	Total
Max	0.1	0.1	0.1	0.1	0.2
Nash	0.1	0.1	0.1	0.1	0.2
Nash + IVH	0.1	0.1	0.1	0.1	0.2

### Sensitivity analysis: Cycle



	$z_1$	$z_2$	$z_3$	$z_4$	Total
Max	0.1	0.1	0.1	0.1	0.2
Nash	0.1	0.1	0.1	0.1	0.2
Nash + IVH	0.1	0.1	0.1	0.1	0.2

Change  $\pi_1^* = 0.2$ 

	$z_1$	$z_2$	$z_3$	$z_4$	Total
Max	0.1533	0.1067	0.1133	0.1267	0.25
Nash	0.1533	0.1067	0.1133	0.1267	0.25
Nash + IVH	0.1644	0.1000	0.1000	0.1178	0.2411

### Sensitivity analysis: Cycle



	$z_1$	$z_2$	$z_3$	$z_4$	Total
Max	0.1	0.1	0.1	0.1	0.2
Nash	0.1	0.1	0.1	0.1	0.2
Nash + IVH	0.1	0.1	0.1	0.1	0.2

Set  $\pi_3^* = 0$  and  $P^{(1)}(3,4) = P^{(0)}(3,4) = 0.5$ 

	$z_1$	$z_2$	$z_3$	$z_4$	Total
Max	0.0867	0.0733	0.0000	0.1400	0.1500
Nash	0.0867	0.0733	0.0467	0.0933	0.1500
Nash + IVH	0.0860	0.0908	0.0000	0.0527	0.1147

# **Concluding remarks**

Principled framework for attribution

- cost approach
- build on accounting principles
- captures conversion funnel and network effect

Current status

- obtain set of feasible costs per impressions
- avoids limitations of last touch and incremental value heuristics

Further questions

- Alternative criteria for identifying unique  $\boldsymbol{z}$
- Empirical tests



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